

# Greedy optimal control for elliptic equations

Applications to turnpike control

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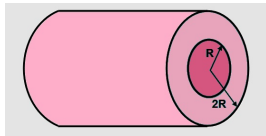
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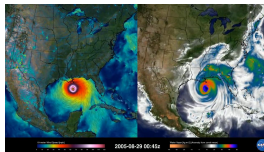
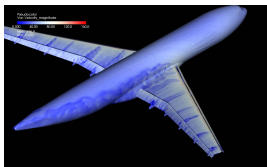


# Parameter dependent problems

Real life applications (may) depend on a **large** number of parameters



**examples:** thickness, conductivity, density, length, humidity, pressure, curvature, . . .



# Parameter dependent problems (Cont.)

- When dealing with applications and simulations, we would like to explore within different parameter configurations.
- From the **control point of view**, this implies solving a different problem **for each configuration**.
- Computationally expensive.

## OUR GOAL

Apply greedy theory to have a robust and fast numerical solvers.

# Parameter dependent control problem

$$\Omega \subset \mathbb{R}^N, \quad \omega \subset \Omega.$$

Consider the system

$$\begin{cases} -\operatorname{div}(a(x, \nu) \nabla y) + c y = \chi_{\omega} u & \text{in } \Omega, \\ y = 0 & \text{on } \partial\Omega, \end{cases} \quad (1)$$

○  $\nu$  is a parameter    ○  $u \in L^2(\omega)$  is a control    ○  $c = c(x) \in L^\infty(\Omega)$

## Optimal control problem ( $\text{OCP}_\nu$ )

$$\min_{u \in L^2(\omega)} J_\nu(u) = \frac{1}{2} |u|_{L^2(\omega)}^2 + \frac{\beta}{2} \|y - y_d\|_{L^2(\Omega)}^2,$$

# Parameter dependent control problem (cont.)

## Optimal control problem (OCP <sub>$\nu$</sub> )

$$\min_{u \in L^2(\omega)} J_{\nu}(u) = \frac{1}{2} |u|_{L^2(\omega)}^2 + \frac{\beta}{2} \|y - y_d\|_{L^2(\Omega)}^2,$$

∃! optimal solution is well-known (Lions, Tröltzsch, . . .)

**Characterization:** optimal pair  $(\bar{u}, \bar{y})$

$$\bar{u} = -\chi_{\omega} \bar{q}$$

$$\begin{cases} -\operatorname{div}(a(x, \nu) \nabla \bar{y}) + c \bar{y} = -\chi_{\omega} \bar{q}, & \text{in } \Omega, \\ -\operatorname{div}(a(x, \nu) \nabla \bar{q}) + c \bar{q} = \beta (\bar{y} - y^d), & \text{in } \Omega, \\ \bar{y} = \bar{q} = 0, & \text{on } \partial\Omega. \end{cases} \quad (2)$$

As the state  $y$  depends on  $\nu$ , also the control  $u$  depends on  $\nu$ .

# Parameter dependent control problem (cont.)

$$\begin{cases} -\operatorname{div}(a(x, \nu) \nabla y) + c y = \chi_{\omega} u & \text{in } \Omega, \\ y = 0 & \text{on } \partial\Omega, \end{cases}$$

From the practical point of view,

- Measure parameter  $\nu$  and determine  $u_{\nu}$

$$\min_{u \in L^2(\omega)} J_{\nu}(u) = \frac{1}{2} \|u\|_{L^2(\omega)}^2 + \frac{\beta}{2} \|y - y_d\|_{L^2(\Omega)}^2,$$

using classical methods (iterative methods, ...)

- Repeat the process for each new value of  $\nu$ .

## CAN WE DO IT BETTER?

# Greedy control

Assume that  $\nu$  ranges within a compact set  $\mathcal{K} \subset \mathbb{R}^d$  and  $a_\nu = a(x, \nu)$  are bounded functions satisfying

$$0 < a_1 \leq a_\nu \leq a_2, \quad \nu \in \mathcal{K}.$$

In this way, we ensure that each control can be uniquely determined by

$$\bar{u}_\nu = -\chi_\omega \bar{q}$$

where  $(\bar{y}, \bar{q})$  solve the optimality system (14). Consider the set of controls  $\bar{u}_\nu$  for each possible value  $\nu \in \mathcal{K}$ . That is,

$$\bar{\mathcal{U}} = \{\bar{u}_\nu : \nu \in \mathcal{K}\}$$




## THE IDEA

To determine a finite number of values of  $\nu$  that yield the best possible approximation of the control manifold  $\bar{\mathcal{U}}$

# Description of the method

We look for a *small* number of parameters  $\nu \in \mathcal{K}$  approximating the manifold  $\bar{\mathcal{U}}$  in the sense of the **Kolmogorov width**. **Roughly**, the **Kolmogorov width** measures how well we can approximate  $\bar{\mathcal{U}}$  by a finite dimensional space.

In order to achieve this goal we rely on **greedy algorithms** and **reduced bases methods** for parameter dependent PDEs or abstract equations in Banach spaces.

-  A. COHEN, R. DEVORE, Kolmogorov widths under holomorphic mappings, *IMA Journal on Numerical Analysis*, to appear
-  A. COHEN, R. DEVORE, Approximation of high-dimensional parametric PDEs, arXiv preprint, 2015.
-  Y. MADAY, O. MULA, A. T. PATERA, M. YANO, The generalized Empirical Interpolation Method: stability theory on Hilbert spaces with an application to the Stokes equation, submitted



# The pure greedy method

$X$  – a Banach space  $K \subset X$  – a compact subset.

The method approximates  $K$  by a series of finite dimensional linear spaces  $V_n$  (a **linear method**).

## THE ALGORITHM

**The first step** Choose  $x_1 \in K$  such that

$$\|x_1\|_X = \max_{x \in K} \|x\|_X.$$

**The general step** Having found  $x_1 \dots x_n$ , denote  $V_n = \text{span}\{x_1, \dots, x_n\}$ .  
Choose the next element

$$x_{n+1} := \arg \max_{x \in K} \text{dist}(x, V_n). \quad (3)$$

**The algorithm stops** when  $\sigma_n(K) := \max_{x \in K} \text{dist}(x, V_n)$  becomes less than the given tolerance  $\varepsilon$ .

# The greedy idea

# The greedy idea

Which one you are going to choose?



Sometimes it is hard to solve the maximisation problem (3).

# The weak greedy method

- a relaxed version of the pure one.

## THE ALGORITHM

Fix a constant  $\gamma \in \langle 0, 1 \rangle$ .

The first step Choose  $x_1 \in K$  such that

$$\|x_1\|_X \geq \gamma \max_{x \in K} \|x\|_X.$$

The general step

Having found  $x_1 \dots x_n$ , denote  $V_n = \text{span}\{x_1, \dots, x_n\}$ .

Choose the next element

$$\text{dist}(x_{n+1}, V_n) \geq \gamma \max_{x \in K} \text{dist}(x, V_n). \quad (4)$$

The algorithm stops when  $\sigma_n(K) := \max_{x \in K} \text{dist}(x, V_n)$  becomes less than the given tolerance  $\varepsilon$ .

In order to estimate **the efficiency of the (weak) greedy algorithm** we compare its approximation rates  $\sigma_n(K)$  with the best possible one.

## The Kolmogorov $n$ width, $d_n(K)$

– measures how well  $K$  can be approximated by a subspace in  $X$  of a fixed dimension  $n$ .

$$d_n(K) := \inf_{\dim Y = n} \sup_{x \in K} \inf_{y \in Y} \|x - y\|_X.$$

Thus  $d_n(K)$  represents optimal approximation performance that can be obtained by a  $n$ -dimensional linear space.

The greedy approximation rates have same decay as the Kolmogorov widths.

- The set  $K$  in general consists of infinitely many vectors.
- In practical implementations the set  $K$  is often unknown (e.g. it represents the family of solutions to parameter dependent problems).

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PRACTICAL REALISATION DEPENDS CRUCIALLY ON AN  
EXISTENCE OF AN APPROPRIATE SURROGATE .

The vectors chosen by the greedy procedure are the **snapshots**.

Their computation can be time consuming and computationally expensive (**offline part**).



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Once having chosen the snapshots, one should easily approximate any value  $x \in K$  (**online part**).

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# The surrogate

In practical implementations, the set  $\bar{\mathcal{U}}$  is **unknown**.

Given two parameters  $\nu_1$  and  $\nu_2$ , how can we measure the distance between  $\bar{u}_{\nu_1}$  and  $\bar{u}_{\nu_2}$ ?

Recall that we want to avoid to compute  $\bar{u}_{\nu}$ .

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**Standard residual:** Suppose that we have computed  $u_{\nu_1}$

$$|u_{\nu_1} - u_{\nu_2}| \sim |\nabla J_{\nu_2}(u_{\nu_1}) - \nabla J_{\nu_2}(u_{\nu_2})| = \nabla J_{\nu_2}(u_{\nu_1})$$

Compute  $\nabla J_{\nu_2}(u_{\nu_1}) = u_{\nu_1} + \beta S_{\nu_2}^*(S_{\nu_2} u_{\nu_1} - y_d)$ , where  $S_{\nu}$  is to control-to-state operator. This means

$$\begin{cases} -\operatorname{div}(a_{\nu_2} \nabla y) + c y = \chi_{\omega} u_{\nu_1}, & \text{in } \Omega, \\ -\operatorname{div}(a_{\nu_2} \nabla q) + c q = \beta (y - y^d), & \text{in } \Omega, \\ y = q = 0, & \text{on } \partial\Omega. \end{cases} \quad \Rightarrow -\chi_{\omega} q = S_{\nu_2}^*(S_{\nu_2} u_{\nu_1} - y_d)$$

# Cheaper surrogates

**A cheap surrogate:** Instead of using  $\bar{u}_\nu$  and approximate the manifold  $\bar{\mathcal{U}}$ , use the optimal variables  $(\bar{q}_\nu, \bar{y}_\nu)$  and approximate the manifold  $\bar{\mathcal{Q}} \times \bar{\mathcal{Y}}$ .

Denoting  $L_\nu z := -\operatorname{div}(a_\nu \nabla z) + c z$ , we define

$$R_\nu(q, y) := \begin{pmatrix} L_\nu y + \chi_\omega q \\ L_\nu q - \beta(y - y_d) \end{pmatrix} = \underbrace{G_\nu(q, y)}_{\text{linear part of } R_\nu(q, y)} + \begin{pmatrix} 0 \\ \beta y_d \end{pmatrix}.$$

With this definition, we are able to compute the following estimates:

$$c_1 \left( \|y - \bar{y}_\nu\|_{H_0^1 \Omega} + \|q - \bar{q}_\nu\|_{H_0^1(\Omega)} \right) \leq \|R_\nu(q, y)\|_{H^{-1}(\Omega)},$$

$$\|R_\nu(p, y)\|_{H^{-1} \Omega} \leq (1 + \alpha_2) (\|y - \bar{y}_\nu\|_{H_0^1(\Omega)} + \|q - \bar{q}_\nu\|_{H_0^1(\Omega)}).$$

where  $c_1$  and  $\alpha_2$  only depending on  $a_1$ ,  $a_2$  and  $\|c\|_\infty$ .

Upper and lower bounds for  $R_\nu(q, y)$  are essential for the proof of greedy algorithms in terms of the Kolmogorov width.

$$R_\nu(q, y) := \begin{pmatrix} L_\nu y + \chi_\omega q \\ L_\nu q - \beta(y - y_d) \end{pmatrix}. \quad (5)$$

## Theorem 1 (H. Santmaria, L., Zuazua, '17)

The residual (5) provides the approximation estimates for optimal controls and states

- $\|u_\nu^* - \bar{u}_\nu\|_{L^2(\Omega)} \leq \frac{1}{c_1} \|R_\nu(q, y)\|_{[H^{-1}(\Omega)]^2},$
- $\|y_\nu^* - \bar{y}_\nu\|_{H_0^1(\Omega)} \leq \left( \frac{1}{\alpha_1 c_1} \right) \|R_\nu(q, y)\|_{[H^{-1}(\Omega)]^2},$

where  $c_1$  and  $\alpha_1$  only depend on  $a_1$ ,  $a_2$  and  $\|c\|_\infty$ .



# Offline algorithm

**Step 1: Initialization.** Fix  $\varepsilon > 0$ . Choose any  $\nu \in \mathcal{K}$ ,  $\nu = \nu_1$  and compute the minimizer of  $J_{\nu_1}$ . This leads to  $\begin{pmatrix} \bar{q}_{\nu_1} \\ \bar{y}_{\nu_1} \end{pmatrix}$ .

**Step 2: recursive choice of  $\nu$ .**

Assuming we have chosen  $\nu_1, \dots, \nu_p$ , we choose  $\nu_{p+1}$  as the maximizer of

$$\max_{\nu \in \mathcal{K}} \left\| \inf_{(q,y) \in (\bar{\mathcal{Q}}_p, \bar{\mathcal{Y}}_p)} R_\nu(q, y) \right\|$$

**Step 3: Stopping criterion.** Stop if the  $\max \leq \varepsilon$ .

## Theorem 1 ( H-Santamaria, L., Zuazua, '17)

The offline algorithm stops after  $n_0(\varepsilon)$  iterations, and fullfills the requirements of the greedy theory.

After choosing the most representative values of  $\nu$ , we can construct an approximated optimal control  $u_\nu^*$  for any arbitrary given value  $\nu \in \mathcal{K}$  by taking

$$u_{\nu}^* = \sum_{i=1}^k \lambda_i \bar{q}_{\nu_i} |_{\omega}$$

where  $\lambda_i$  are determined by the projection of the vector  $\begin{pmatrix} 0 \\ y_d \end{pmatrix}$  to the space

$$\text{span}\{G_\nu(\bar{q}_{\nu_1}, \bar{y}_{\nu_1}), \dots, G_\nu(\bar{q}_{\nu_k}, \bar{y}_{\nu_k})\}$$

# NUMERICAL RESULTS

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# Numerical examples

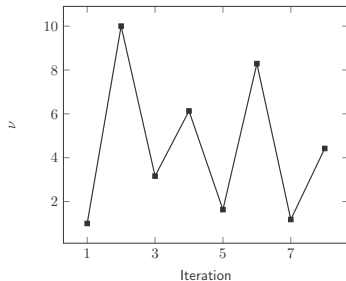
- $\Omega = (0, 1)^2$  in 2-D or  $\Omega = (0, 1)$  in 1-D.
- Uniform meshes, i.e., meshes with constant discretization steps in each direction,  $N = 400$ .
- We will approximate the operator  $\mathcal{A} = -\operatorname{div}(a(x, \nu)\nabla \cdot)$  by using the standard 5-point discretization.
- Discretize-then-optimize.
- $\nu \in \mathcal{K} = [1, 10]$ .
- $\mathcal{K}$  sampled in 100 equidistant points.

# Greedy test # 1

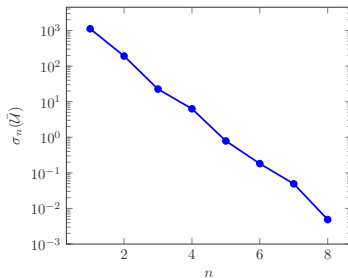
$$\circ a(x, \nu) = 1 + \nu(x_1^2 + x_2^2), \quad \circ c(x) = \sin(2\pi x_1) \sin(2\pi x_2),$$

$$\circ y_d = \sin(\pi x_1), \quad \circ \beta = 10^4, \quad \circ \varepsilon = 0.005$$

$$\circ t_{\text{cheap}} = 304s, \quad \circ t_{\text{std}} = 384s$$

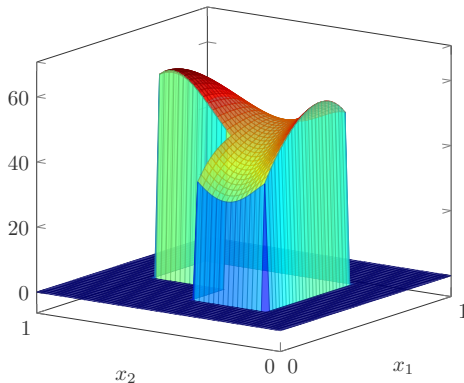


(a) Selected  $\nu$

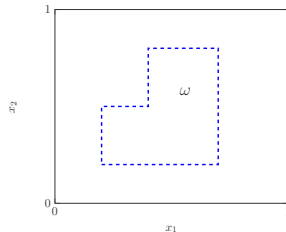


(b) Approximation error

# Approximation for $\nu = \pi/2$



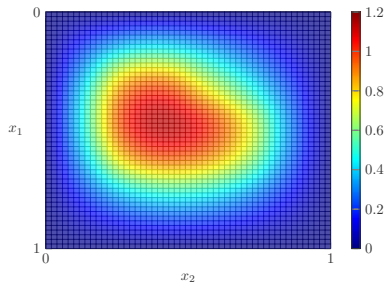
(c) The approximated control



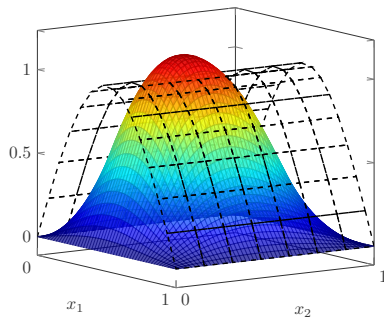
(d) The control set  $\omega$

$$\circ |u_{\pi/2}^* - \bar{u}_{\pi/2}|_{L^2(\omega)} \approx 1.45 \times 10^{-5}, \quad t_{\text{online}} = 0.45s, \quad t_{\text{iterative}} = 6.01s.$$

# Approximation for $\nu = \pi/2$ (cont.)



(e) The controlled state



(f) The state  $y_{\pi/2}^*$  and the target function  $y^d$  (dashed)

$$\circ |y_{\pi/2}^* - \bar{y}_{\pi/2}|_{L^2(\Omega)} \approx 1.15 \times 10^{-7}$$

# CONNECTION WITH THE TURNPIKE PROBLEMS

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# Time dependent control problem

Consider

$$\begin{cases} \partial_t y - \operatorname{div}(a(x, \nu) \nabla y) + c y = \chi_\omega u & \text{in } Q = \Omega \times (0, T), \\ y = 0 & \text{on } \Sigma = \partial\Omega \times (0, T), \\ y(x, 0) = y^0(x) & \text{in } \Omega. \end{cases} \quad (6)$$

and the control problem

$$\min_u J_\nu^T(u) = \frac{1}{2} \int_0^T |u(t)|_{L^2(\omega)}^2 dt + \frac{\beta}{2} \int_0^T \|y(t) - y^d\|_{L^2(\Omega)}^2 dt.$$

The optimal solution  $(u^T, y^T)$  satisfies

$$\|y^T(t) - \bar{y}\|_{L^2(\Omega)} + \|u^T(t) - \bar{u}\|_{L^2(\Omega)} \leq K \left( e^{-\mu t} + e^{-\mu(T-t)} \right), \quad \forall t \in [0, T]$$

- Exponential convergence of the finite-time horizon control problem to the steady one as  $T \rightarrow \infty$ .

## Greedy test # 2

We consider time-dependent version of the last example.

$$\circ a(x, \nu) = 1 + \nu(x_1^2 + x_2^2), \quad \circ c(x) = \sin(2\pi x_1) \sin(2\pi x_2),$$

$$\circ y_d = \sin(\pi x_1), \quad \circ \beta = 10^4, \quad \circ \varepsilon = 0.005$$

$$\circ \Omega = (0, 1)^2 \text{ in 2-D or } \Omega = (0, 1) \text{ in 1-D.}$$

We take initial datum

$$y_0(x) = \sin(3\pi x_1) \sin(2\pi x_2)$$

## The case $c(x) \geq 0$ (greedy test #2)

$$u(x, t) = u_{\pi/2}^*(x), \quad y_0(x) = \sin(3\pi x_1) \sin(2\pi x_2)$$

